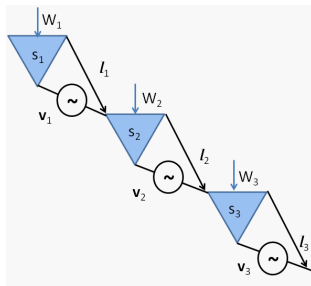


# Mixed Integer Dynamic Approximation Scheme

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(Joint work with Faisal Wahid, Frederic Bonnans)

# Motivation: daily hydro generation in a river chain



Given hourly electricity prices  $\pi(t)$  the generator arranges releases  $v_i(t)$  and spill  $l_i(t)$  of water to maximize revenue  $\sum_t \pi(t) \sum_i g_i(v_i(t))$  while respecting the water flow constraints of the river chain. Here  $g_i$  converts water flow into power.

# Autoregressive price model and random inflows

Given initial state  $\mathbf{x}_0$ , we seek an optimal policy yielding  $V_1(\mathbf{x}_0)$ , where

$$V_t(\mathbf{x}) = \mathbb{E}_{\xi_t} \left[ \max_{\mathbf{u} \in U(\mathbf{x})} \{ r_t(\mathbf{x}, \mathbf{u}, \xi_t) + V_{t+1}(f_t(\mathbf{x}, \mathbf{u}, \xi_t)) \} \right],$$

$$V_{T+1}(\mathbf{x}) = R(\mathbf{x})$$

$$\mathbf{x} = [ s \quad \pi ]^\top, \quad \mathbf{u} = [ v \quad l ]^\top, \quad \xi_t = [ \omega_t \quad \eta_t ]^\top$$

$$f_t(s, \pi, v, l, \omega_t, \eta_t) = \begin{bmatrix} s_t - v_t - l_t + \omega_t \\ \alpha_t \pi_t + (1 - \alpha_t) b_t + \eta_t \end{bmatrix},$$

$$r_t(s, \pi, v, l, \omega_t, \eta_t) = \pi \sum_i g_i(v).$$

# How to solve this

- Stochastic programming (Fleten and Kristoffersen, 2007)
- Backward recursion (Pritchard and Zakeri, 2003)
- SDDP (Pereira and Pinto, 1991, P., Dallagi, Gallet, 2013)
- ADDP (Löhndorf et al, 2013)
- Linear decision rules (Braathen et al, 2013)
- SDDP with MIPs (Zhou, Ahmed, Sun, 2016)

# Our approach

- Stochastic control using approximate dynamic programming.
- SDDP constructs cutting-plane outer approximation of convex or concave value functions.
- MIDAS: similar methodology to (approximately) solve stochastic optimal control problems with nonconvex value functions (e.g from AR model of price).
- Solves stage problems using mixed integer programming.
- Convergence requires monotonicity and continuity of  $V_t(\mathbf{x})$ .

# Summary

- 1 Introduction
- 2 Outer approximation of value function
- 3 MIDAS
- 4 Computational results
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# Dynamic programming formulation

Given initial state  $x_0$ , we seek an optimal policy yielding  $V_1(x_0)$ , where

$$V_t(x) = \mathbb{E}_{\zeta_t} \left[ \max_{u \in U(x)} \{r_t(x, u, \zeta_t) + V_{t+1}(f_t(x, u, \zeta_t))\} \right]$$
$$V_{T+1}(x) = R(x).$$

Here  $V_t(x)$  denotes the maximum expected reward from the beginning of stage  $t$  onwards, given the state is  $x$ , and we take action  $u_t$  after observing the random disturbance  $\zeta_t$ . We assume that  $R(x)$  is continuous, and  $U(x)$  is sufficiently regular so that  $V_t$  is continuous if  $V_{t+1}$  is.

Given  $\varepsilon > 0$ , there is some  $\delta$  so that for all  $t = 1, 2, \dots, T + 1$ ,

$$\|x - y\|_\infty < \delta \Rightarrow |V_t(x) - V_t(y)| < \varepsilon.$$



# Outer approximation of continuous monotonic functions

Given a continuous nondecreasing function  $Q(x) \leq M$ , and a finite set of values

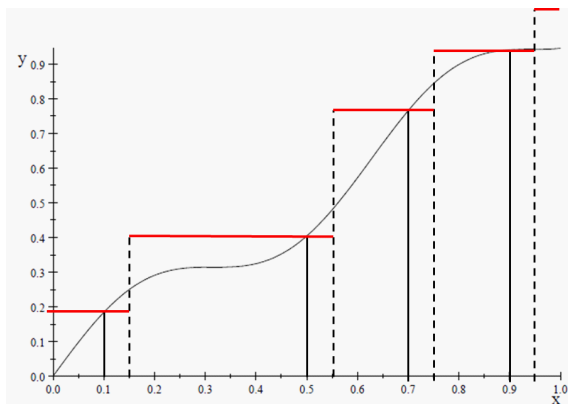
$$Q(x^h) = q^h, \quad h = 1, 2, \dots, H,$$

approximate  $Q(x)$  by a piecewise constant function  $Q^H(x)$  so that for every  $x$

$$Q(x) \leq Q^H(x) + \varepsilon$$

- $q^h$  is a real number and  $Q^H$  is a function;
- $Q(x)$  is assumed monotonic to guarantee that  $Q(x) \leq Q^H(x) + \varepsilon$  for every  $x$ .

# Example



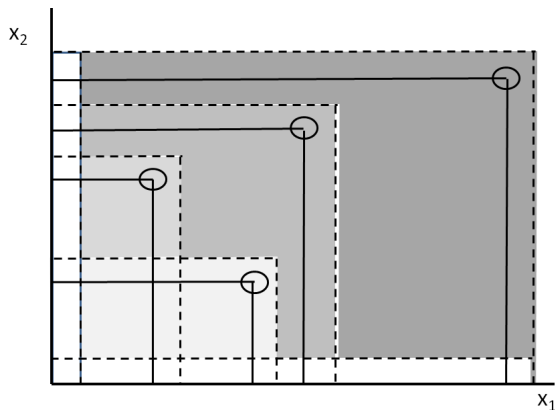
Approximation of  $Q(x) = x + 0.1 \sin(10x)$  at points  $x^h = 0.1, 0.5, 0.7, 0.9$ , and  $\delta = 0.05$ .  $Q^H(x)$  shown in red is upper semicontinuous, and is an upper bound on  $Q(x) - \varepsilon$ .

# MIP approximates a continuous monotonic function

Assume that

$$\begin{array}{ll}
 Q^H(x) = \max & \varphi \\
 \text{s.t.} & \varphi \leq q^h + M(1 - y^h), \quad h = 1, 2, \dots, H, \\
 & x_k \geq x_k^h z_k^h + \delta, \quad k = 1, 2, \dots, n, \\
 & \phantom{x_k \geq} \phantom{x_k^h z_k^h + \delta}, \quad h = 1, 2, \dots, H, \\
 & \sum_{k=1}^n z_k^h = 1 - y^h, \quad h = 1, 2, \dots, H, \\
 & y^h \in \{0, 1\}, \quad h = 1, 2, \dots, H, \\
 & z_k^h \in \{0, 1\}, \quad k = 1, 2, \dots, n, \\
 & \phantom{z_k^h \in} \phantom{z_k^h \in}, \quad h = 1, 2, \dots, H.
 \end{array}$$

# Example in two dimensions

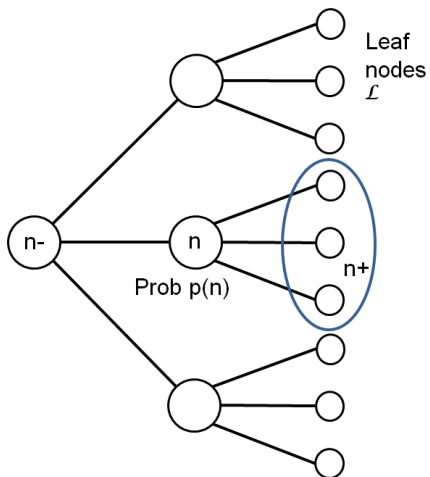


Contour plot of  $Q^H(x)$  when  $H = 4$ . Circled points are  $x^h$ ,  $h = 1, 2, 3, 4$ . Darker shading indicates increasing values of  $Q^H(x)$  that equals  $Q(x^h)$  in each region containing  $x^h$ ,  $h = 1, 2, 3, 4$ .  $Q^H(x)$  defined only when  $x \geq \delta e$ .

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# Multistage algorithm uses a scenario tree



# Tree formulation of stochastic control problem

$$\begin{aligned}
 \text{MSPT: } \max \quad & \sum_{n \in \mathcal{N} \setminus \{0\}} p(n) r_n(x_{n-}, u_n) + \sum_{n \in \mathcal{L}} p(n) R(x_n) \\
 \text{s.t. } \quad & x_n = f_{n-}(x_{n-}, u_n, \xi_n), \\
 & x_0 = \bar{x}, \\
 & u_n \in U(x_n), \\
 & x_n \in X_n.
 \end{aligned}$$

DP recursion is:

$$\begin{aligned}
 V_n(x_n) &= \sum_{m \in n^+} \frac{p(m)}{p(n)} \max_{u \in U(x_n)} \{r_m(x_n, u) + V_m(f_n(x_n, u, \xi_m))\} \\
 V_n(x_n) &= R(x_n), \quad n \in \mathcal{L},
 \end{aligned}$$

where we seek a policy that maximizes  $V_0(x_0)$ .

# MIDAS algorithm

- 1 Set  $Q_n^1(x) = M$ , for every  $n \in \mathcal{N} \setminus \mathcal{L}$ ;
- 2 For  $H = 1, 2, \dots$ ,
  - set  $Q_n^H(x) = R(x)$ , for every  $n \in \mathcal{L}$ ;
  - perform a **forward** pass then a **backward** pass.



# Forward pass

Set  $x_0^H = x_0$ , and  $n = 0$ . While  $n \notin \mathcal{L}$ :

- ① Sample  $m \in n+$  to give  $\zeta_m^H$ ;
- ② Solve  $\max_{u \in U(x_n^H)} \left\{ r_m(x_n^H, u) + Q_m^H(f_n(x_n^H, u, \zeta_m^H)) \right\}$  to give  $u_m^H$ ;
- ③ If  $\left\| f_n(x_n^H, u_m^H, \zeta_m^H) - x_m^h \right\|_\infty < \delta$  for  $h < H$  then set  $x_m^{H+1} = x_m^h$ , else set  $x_m^{H+1} = f_n(x_n^H, u_m^H, \zeta_m^H)$ ;
- ④ Set  $n = m$ .

# Backward pass

For the particular node  $n \in \mathcal{L}$  at the end of forward pass update  $Q_n^H(x)$  to  $Q_n^{H+1}(x)$  by adding  $q_n^{H+1} = R(x_n^{H+1})$  at point  $x_n^{H+1}$ .

While  $n > 0$

- 1 Set  $n = n -$ ;
- 2 Compute

$$\varphi_n = \sum_{m \in n^+} \frac{p(m)}{p(n)} \max_{u \in U(x_n^H)} \left\{ r_m(x_n^{H+1}, u) + Q_m^{H+1}(f_n(x_n^{H+1}, u, \zeta_m)) \right\}$$

- 3 Update  $Q_n^H(x)$  to  $Q_n^{H+1}(x)$  by adding  $q_n^{H+1} = \varphi_n$  at point  $x_n^{H+1}$ ;

# Sampling property

**FPSP:** For each  $n \in \mathcal{L}$ , with probability 1

$$\left| \left\{ H : \zeta_n^H = \zeta_n \right\} \right| = \infty.$$

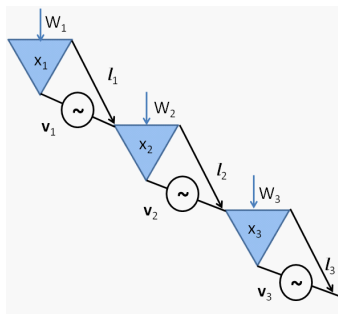
## Theorem

*If step 1 of forward pass satisfies FPSP then sampled MIDAS converges almost surely to a  $(T + 1)\varepsilon$ -optimal solution to MSPT.*

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# Recall hydro generation in a river chain



Given future electricity prices  $\pi(t)$  the generator arranges releases  $v_i(t)$  and spill  $I_i(t)$  of water to maximize revenue  $\sum_t \pi(t) \sum_i g_i(v_i(t))$  while respecting the water flow constraints of the river chain.

## Example: AR1 price model for river chain

$$\mathbf{x}_{t+1} = f_t(\mathbf{x}, \mathbf{u}, \boldsymbol{\zeta}_t),$$

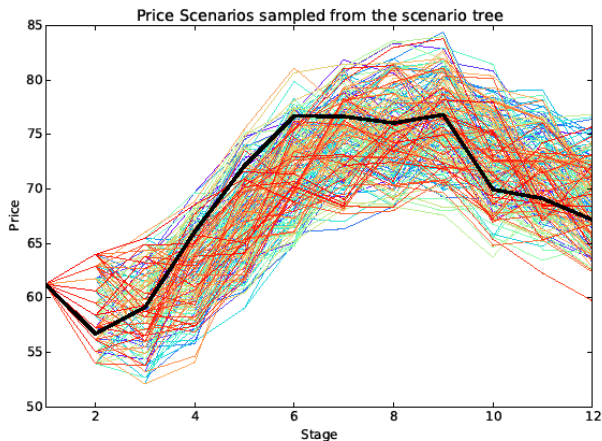
$$\begin{bmatrix} s_{t+1} \\ \pi_{t+1} \end{bmatrix} = \begin{bmatrix} s_t - v_t - l_t + \omega_t \\ \alpha_t \pi_t + (1 - \alpha_t) b_t + \eta_t \end{bmatrix},$$

Here  $\omega_t$  is (random) reservoir inflow,  $\eta_t$  is error term for price model, so  $\boldsymbol{\zeta}_t = [\omega_t \ \eta_t]^\top$  and  $\mathbf{u} = [v \ l]^\top$  release and spill. Reward in stage  $t$  is

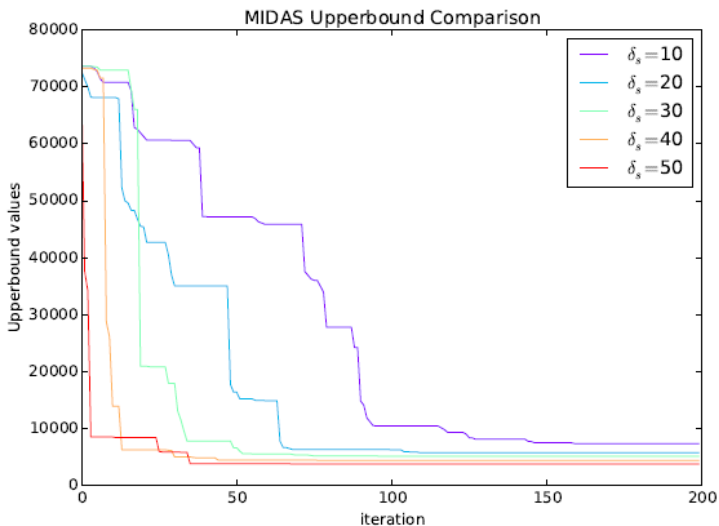
$$r_t(s, \pi, v, l, \omega_t, \eta_t) = \pi \sum_i g_i(v),$$

from released energy  $g(v)$  sold at price  $\pi$ , and  $U(\mathbf{x}) = \{(v, l) : v \in U_0, v + l \in [0, s]\}$ .

# Price scenarios sampled from AR1 model

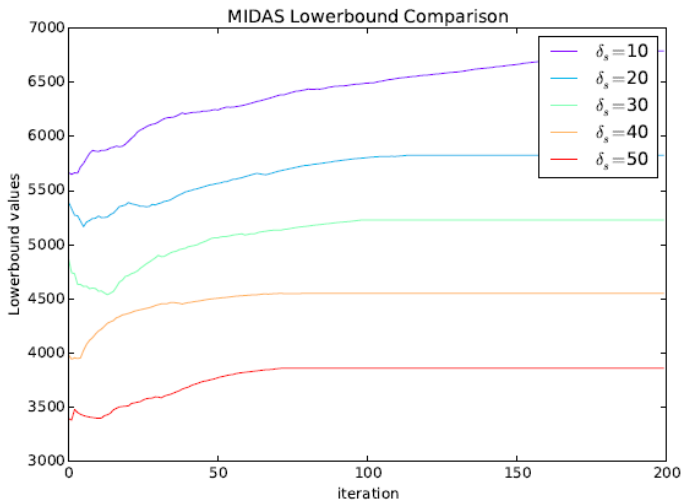


# Epsilon upper bounds for single reservoir problem

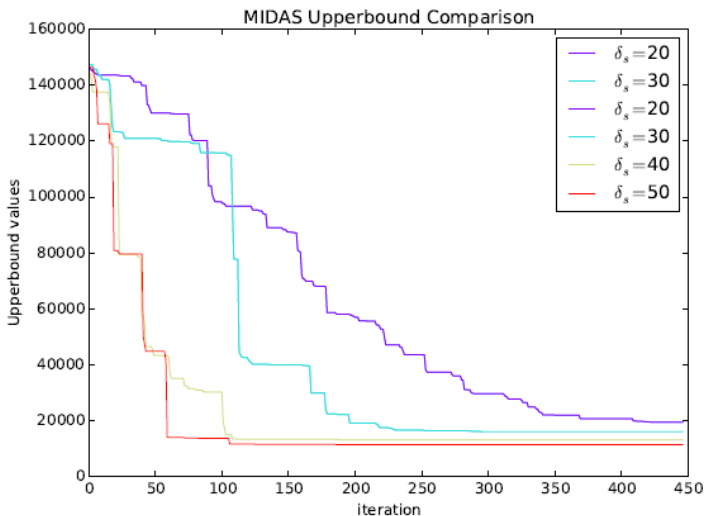




# Estimated policy payoffs for different delta



# Epsilon upper bounds for two-reservoir problem



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# Conclusions

- SDDP has proved very successful in hydrothermal scheduling;
- Sampled trajectories reduce DP computation;
- Concave value functions a limitation;
- MIDAS is an attempt to extend these features to more general stochastic dynamic programs in the hope of making them tractable by solving small MIPs;
- Convergence wp1 can be shown for continuous monotonic value functions.

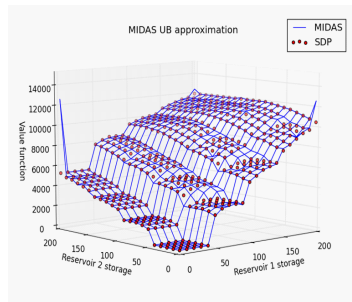
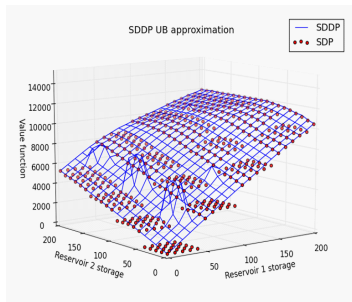
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# SDDP and MIDAS for stochastic MIPs



True value function shown in red. SDDP and MIDAS give different outer approximations when true value function is not concave.

# SDDP and MIDAS

Solved & simulated offer policy of MIDAS & SDDP for range of initial storage levels

**MIDAS policy:** On average 98% of optimal value

**SDDP policy:** On average 93% of optimal value

	<b>Mean</b>	<b>Median</b>	<b>Upper quartile</b>	<b>Lower quartile</b>
<b>SDDP</b>	92.92	88.22	81.18	96.35
<b>MIDAS</b>	97.93	99.06	97.18	99.67



# The one-dimensional case

$$\begin{aligned} V_H(x) = \max \quad & \varphi \\ \text{s.t.} \quad & \varphi \leq V^h + Mz^h, \quad h = 1, 2, \dots, H, \\ & x \geq x^h z^h + \delta, \quad h = 1, 2, \dots, H, \\ & z^h \in \{0, 1\}, \quad h = 1, 2, \dots, H. \end{aligned}$$

Consider  $(x^h, V^h)$ ,  $h = 1, 2, \dots, H$ . The variable  $z^h = 1$  picks out all the  $x^h$  lying at or below  $x - \delta$ . These do not constrain  $\varphi$ . The  $x^h$  lying strictly above  $x - \delta$  have  $z^h = 0$ . Thus

$$V_H(x) = \min\{V^h : x^h > x - \delta\}.$$

# Explanation of approximate stage problem 1

If  $w_h = 0$  for some  $h$  then  $\sum_k z_k^h = 1$ , so  $z_k^h = 1$  for exactly one  $k$ .  
The constraint

$$x_k \geq x_k^h z_k^h + \delta$$

means that there is at least one dimension  $k$  with  $x_k > x_k^h$ . So  $x$  is somewhere in the complement of the rectangular box  $R = \{y \mid y \leq x^h\}$ . Since  $w_h = 0$ , we have

$$\varphi \leq v_h + M$$

so the value  $v^h$  from top right-hand corner of  $R$  does not bound  $\varphi$ .

## Explanation of approximate stage problem 2

If  $x \leq x^h$  for  $h \in \mathcal{G} \subseteq \{1, 2, \dots, H\}$ , then for each  $h \in \mathcal{G}$  we must have every  $z_k^h = 0$ , so  $\sum_k z_k^h = 0$ . It follows from

$$\sum_{k=1}^n z_k^h = 1 - w_h,$$

that  $w_h = 1$ ,  $h \in \mathcal{G}$ , so the constraints become

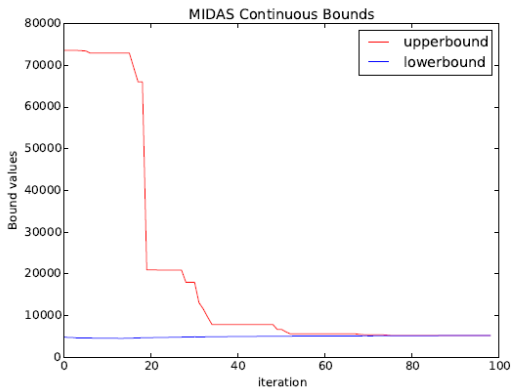
$$\begin{aligned} \varphi &\leq v_h, & h \in \mathcal{G}, \\ x_k &\geq 0 + \delta, & k = 1, \dots, n, \end{aligned}$$

so this will result in

$$\varphi \leq \min\{v_h \mid h \in \mathcal{G}\},$$

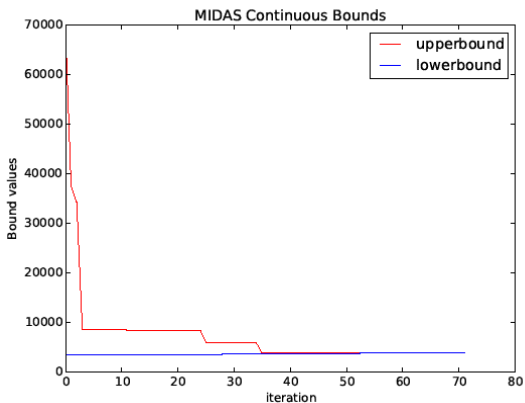
as desired.

# Epsilon upper and lower bounds for single reservoir problem



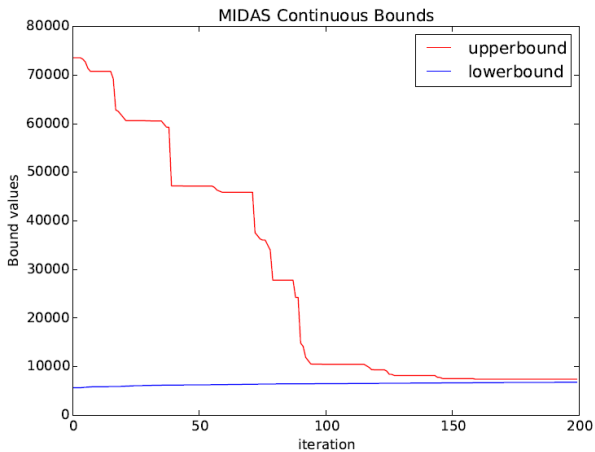
With  $\delta = 30$

# Epsilon upper and lower bounds for single reservoir problem



With  $\delta = 50$

# Epsilon upper and lower bounds for single reservoir problem



With  $\delta = 10$