

Financial transmission rights and SPD

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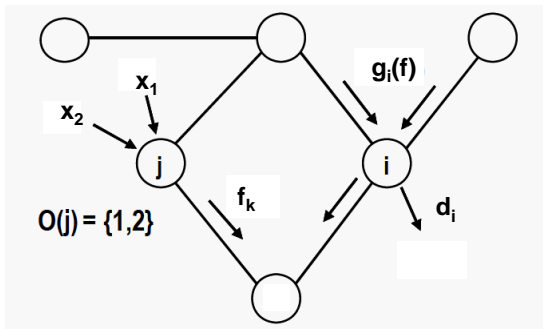
Motivation

- FTR design = economics+politics+mathematics
- How can mathematics help this process?
- What is known mathematically for FTRs in NZ setting?
- Are there any traps?
- Disclaimer: This is a talk about mathematics, not a recommendation for a specific FTR design.

Overview of talk

- An abstract formulation of SPD
- Revenue adequacy of FTR products
- Losses
- Reserve

Classical dispatch model (with losses)



$$\begin{aligned}
 \min \quad & \sum_i \sum_{j \in O(i)} c_j x_j \\
 \text{s.t.} \quad & g_i(f) + \sum_{j \in O(i)} x_j - z_i = d_i, \quad i = 1, 2, \dots, n \quad [\pi_i] \\
 & f \in F, \\
 & x \in X, \\
 & z_i \geq 0, \quad i = 1, 2, \dots, n.
 \end{aligned}$$

$g_i(f)$ = flow from the network into each node (demand)

Lossless flow

$$g_i(f) = \sum_{k \in \mathcal{F}(i)} -f_k + \sum_{k \in \mathcal{T}(i)} f_k$$

Quadratic losses

$$g_i(f) = \sum_{k \in \mathcal{F}(i)} \left(-f_k - \frac{r_k}{2} f_k^2 \right) + \sum_{k \in \mathcal{T}(i)} \left(f_k - \frac{r_k}{2} f_k^2 \right) \quad (1)$$

Piecewise linear losses (SPD)

$$g_i(f) = \sum_{k \in \mathcal{F}(i)} -f_k + \sum_{k \in \mathcal{T}(i)} (f_k - T(f_k))$$

What is the set F ?

In the classical case

$$f \in F$$

amounts to

$$-K \leq f \leq K, \quad \text{thermal limits}$$

$$Uf = 0, \quad \text{loop flow constraints (or similar)}$$

We might add

$$Wf \leq a, \quad \text{branch security constraints}$$

But we do not allow constraints mixing f with other variables e.g.

$$\sum_i \sum_{j \in O(i)} x_j \geq s^T f$$

Rentals

Suppose that the optimal dispatch is x^* . Generators of energy earn a *generator rental*

$$R_g = \sum_i \pi_i \sum_{j \in O(i)} x_j^* - \sum_i \sum_{j \in O(i)} c_j x_j^*.$$

The difference between what is paid to the system clearing manager by loads and what they pay generators is the *transmission rental*

$$R_t = \sum_i \pi_i d_i - \sum_i \pi_i \sum_{j \in O(i)} x_j^*.$$

Transmission rental formula

For an optimal feasible dispatch

$$g_i(f^*) + \sum_{j \in O(i)} x_j^* - z_i^* = d_i$$

$$\begin{aligned} R_t &= \sum_i \pi_i d_i - \sum_i \pi_i \sum_{j \in O(i)} x_j^* \\ &= \sum_i \pi_i (g_i(f^*) - z_i^*) \\ &= \sum_i \pi_i g_i(f^*) \end{aligned}$$

Revenue inequality theorem

For any feasible set of flows ($f \in F$)

$$\sum_i \pi_i g_i(f^*) \geq \sum_i \pi_i g_i(f).$$

The transmission rental R_t collected from the optimal dispatch equals or exceeds $\sum_i \pi_i g_i(f)$ for any set of branch flows in F .

Financial Transmission Rights (FTRs)

What is a *financial transmission right* (mathematically)?

It is a vector $h(\alpha)$ of "offtakes" at the nodes of the network, where $h_i(\alpha)$ is the offtake at node i . The FTR $h(\alpha)$ generates a payment to its holder of

$$\sum_i \pi_i h_i(\alpha)$$

For simplicity, we will restrict attention to FTRs of the form

$$h(\alpha) = [-r_1 \quad 0 \quad \dots \quad 0 \quad r_n]^\top$$

A *balanced* FTR $h(\alpha)$ has $\sum_i h_i(\alpha) = 0$.

Simultaneous feasibility

Suppose we allocate some FTRs as vectors $h(\alpha)$, $\alpha \in A$. Then these are *simultaneously feasible* if

$$\begin{aligned} \text{SF: } g_i(v) - z_i &= \sum_{\alpha} h_i(\alpha), & i = 1, 2, \dots, n, \\ z_i &\geq 0, & i = 1, 2, \dots, n, \\ v &\in F, \end{aligned}$$

has a feasible solution v, z .

Observations

- A set of FTRs can be defined by a flow vector v
- In SPD framework, v is by convention the flow *leaving* each node.
- The vector v can be anything in F .
- In SPD we have

$$g_i(v) = \sum_{k \in \mathcal{F}(i)} -v_k + \sum_{k \in \mathcal{T}(i)} (v_k - T(v_k))$$

where $T(v)$ are piecewise linear transmission losses.

- The total payout for the set of FTRs is $\sum_i \pi_i g_i(v)$.

Simultaneous feasibility implies revenue adequacy

- Consider FTRs between two nodes B and O with prices π_B and π_O .
- Suppose a set of FTRs is now defined by $\{(-h_B(\alpha), h_O(\alpha)), \alpha \in A\}$
- Aggregate to get $(-\sum_{\alpha} h_B(\alpha), \sum_{\alpha} h_O(\alpha))$
- Find a flow $v \in F$ that satisfies SF.
- This implies

$$\pi_O \sum_{\alpha} h_O(\alpha) - \pi_B \sum_{\alpha} h_B(\alpha) = \sum_i \pi_i g_i(v) \leq \sum_i \pi_i g_i(f^*)$$

Example

Price	5					OfferQ	Offer P	
						50	5	
						22	110	
						22		
	Tranche Flow	Tranche Capacity	Loss Slope					
	5	5	0.1	4.5	4	Line	Rentals	50
	5	5	0.2	4	3			
	5	5	0.3	3.5	2		Received	160
	5	5	0.4	3	1		Paid	110
	2	5	0.5	1	0		Check	50
Sent		22	Received	16				
						Loss =		6
Price	10	Demand	Supply					
		16	16					

An FTR of $(-20, 15)$ is simultaneously feasible. It earns a payment of $\$10 \cdot 15 - \$5 \cdot 20 = \$50$.

An FTR of $(-20, 14)$ is also simultaneously feasible. It earns a payment of $\$10 \cdot 14 - \$5 \cdot 20 = \$40$.

Simultaneous feasibility with losses

Any set of balanced FTRs will not be simultaneously feasible.

$$\begin{aligned} \text{SF: } g_i(v) - z_i &= -r, & i = 1, \\ g_i(v) - z_i &= 0, & i = 2, \dots, n-1, \\ g_i(v) - z_i &= r, & i = n, \\ z_i &\geq 0, & i = 1, 2, \dots, n, \\ v &\in F \end{aligned}$$

is feasible only if

$$\sum_{i=1}^n g_i(v) \geq 0$$

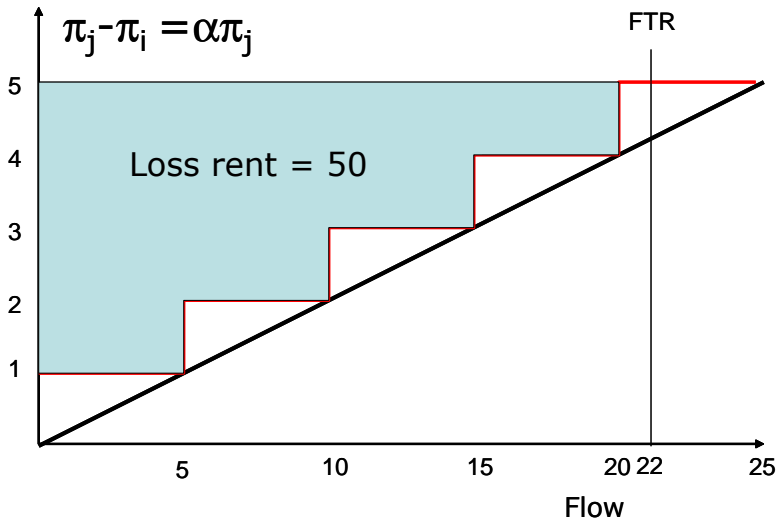
which cannot happen if there are any nonzero losses.

Revenue inadequacy with losses

Price	5					OfferQ	Offer P		
						50	5		
						22	110		
						22			
	Tranche	Tranche	Loss						
	Flow	Capacity	Slope						
	5	5	0.1	4.5	4	Line	Rentals	50	
	5	5	0.2	4	3				
	5	5	0.3	3.5	2		Received	160	
	5	5	0.4	3	1		Paid	110	
	2	5	0.5	1	0		Check	50	
Sent		22	Received	16			Loss =	6	
Price	10	Demand	Supply						
		16	16						

A balanced FTR $(-22, 22)$ earns a payment of $\$(10-5)*22=\110 .
 But $R_t = \$50$. The \$60 shortfall is the cost of the 6MWh loss.

In graphical terms



Losses require modifications to classical design

Can do one of the following to ensure revenue adequacy:

- Restriction on volume of FTR allocation.
- Loss support contracts to accompany a balanced FTR.
- Unbalanced FTRs $(-1, 1 - \alpha)$ for some α .
- Discounted balanced FTRs
- Some combination of the above.

Unbalanced and discounted FTRs

An *unbalanced* FTR $(-r, (1 - \alpha)r)$ gives

$$\begin{aligned} \text{SF: } g_i(v) - z_i &= -r, & i = 1, \\ g_i(v) - z_i &= 0, & i = 2, \dots, n-1, \\ g_i(v) - z_i &= (1 - \alpha)r, & i = n, \\ z_i &\geq 0, & i = 1, 2, \dots, n, \\ v &\in F, \end{aligned}$$

which will be feasible for some v and z .

An α -discounted balanced FTR from i to j pays $(1 - \alpha)\pi_j - \pi_i$ per MW of balanced FTR. It is equivalent to an unbalanced FTR.

Reserve-constrained dispatch (SPD)

$$\begin{array}{ll}
 \min & \sum_i \sum_{j \in O(i)} c_j x_j + \sum_i \sum_{j \in R(i)} b_j y_j \\
 \text{s.t.} & g_i(f) + \sum_{j \in O(i)} x_j - z_i = d_i, \quad i = 1, 2, \dots, n \\
 & \sum_i \sum_{j \in R(i)} y_j - w_1 = q^\top x, \\
 & \sum_i \sum_{j \in R(i)} y_j - w_2 = s^\top f, \\
 & f \in F, \\
 & (x, y) \in X, \\
 & z_i \geq 0, \quad i = 1, 2, \dots, n \\
 & w_i \geq 0, \quad i = 1, 2.
 \end{array}$$

This contains constraints mixing f and y and so we cannot apply the SF test as it is.

What is going on here?

- Assume that the funding for reserve is coming from another source (e.g. a levy). So the revenue inadequacy is not about paying for existing reserve from transmission rentals.
- An FTR of up to 20 MW is feasible without interfering with reserve constraint. The payouts are then less than \$300.
- If the FTR exceeds 20MW then it requires more reserve at \$15/MW. This is the shadow price on the reserve constraint. But this extra money does not come from the dispatch.
- Some reserve support is needed if we want to allocate more than 20MW of FTR
- Observe that it does not help if we consider the South generator as the reserve provider.

Conclusions

- Classical results ignore losses and ancillary constraints - mathematics is then "simple".
- Balanced FTRs are less complex for trading, so more liquid.
- Disadvantage is the complexity in ensuring that they can be funded from rentals.